

Exam Probabilistic Reasoning

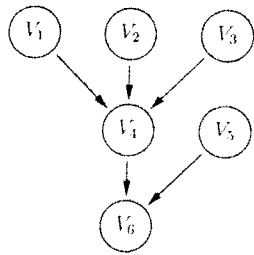
6 November 2013

The exam consists of three problems, for each of which the number of credits per question is given. In total, a maximum of 100 credits is awarded. Read the questions very carefully; you may answer them in English and/or Dutch. Be sure to clearly explain your answers!

Good luck!

Problem 1 (a: 10 pts., b: 10 pts., c: 10 pts., total: 30 pts.)

Consider a probabilistic network $B = (G, \Gamma)$, where $G = (V(G), A(G))$ is the following acyclic digraph and $\Gamma = \{\gamma_{V_i} \mid V_i \in V(G)\}$ is given by:



$$\gamma(v_1) = 0.9$$

$$\gamma(v_2) = 0.4$$

$$\gamma(v_3) = 0.5$$

$$\gamma(v_5) = 0.7$$

$$\gamma(v_4 \mid v_1 \wedge \neg v_2 \wedge \neg v_3) = 0$$

$$\gamma(v_6 \mid v_4 \wedge v_5) = 0.1$$

$$\gamma(v_4 \mid \neg v_1 \wedge v_2 \wedge \neg v_3) = 0.5$$

$$\gamma(v_6 \mid \neg v_4 \wedge v_5) = 0.3$$

$$\gamma(v_4 \mid \neg v_1 \wedge \neg v_2 \wedge v_3) = 0.2$$

$$\gamma(v_6 \mid v_4 \wedge \neg v_5) = 0.7$$

$$\gamma(v_6 \mid \neg v_4 \wedge \neg v_5) = 0.6$$

Variables V_1, V_2 and V_3 have a disjunctive interaction effect on variable V_4 . To capture this effect, the assessment function for node V_4 is based on the ‘noisy-or gate’.

- a. Complete the assessment function $\gamma(V_4)$ for node V_4 . Explain your answers.

Let Pr be the probability distribution defined by probabilistic network B . Now, consider the five computation rules of *Pearl's* data fusion algorithm, given in the attached formula sheet.

- b. Illustrate *Pearl's* algorithm by computing the prior probabilities $\text{Pr}(v_4)$ and $\text{Pr}(\neg v_4)$ from network B .

Clearly indicate which probabilities are computed, which messages are passed and how they are computed; explicitly mention all assumptions you make.

For the next question you may assume that you have an inference algorithm at your disposal (e.g. *Pearl*) which can compute probabilities of the form $\text{Pr}(c_{V_i} \mid c_E)$ ($V_i \in V(G), E \subseteq V(G)$).

- c. You are interested in computing the probability $\text{Pr}(v_1 \wedge \neg v_4 \wedge v_6)$. Clearly explain *how* you can compute this specific probability *efficiently* from network B .

Note: you don't have to perform the actual computation, just explain the approach you would take.

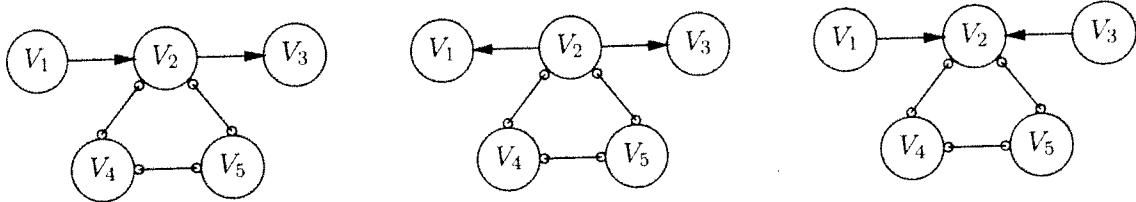
Problem 2 (a: 10 pts., b: 10 pts., c: 15 pts., total: 35 pts.)

Consider a probabilistic network $B = (G, \Gamma)$, a node of interest $T \in V(G)$ and a set of observed nodes $E \subseteq V(G)$ such that $T \notin E$. In a probabilistic network we can identify nodes that are *barren* with respect to T and E : a node $X \in V(G)$ is called *barren* if

$X \neq T$ and $X \notin E$, and all its descendants $\sigma^*(X)$ are barren

- Consider the probabilistic network from **Problem 1**. Let $T = V_2$ and let $E = \emptyset$. Give all nodes that are barren with respect to T and E . Explain your answer.
- To speed up inference, a probabilistic network can be *pruned* to a computationally equivalent network for computing a specific distribution $\Pr(T \mid c_E)$. To this end, all d-separated nodes and all barren nodes with respect to T and E are removed. Clearly explain why these two sets of nodes can be safely removed for computing $\Pr(T \mid c_E)$.

Consider the following graphical structures, where each \circ indicates the presence or absence of an arrowhead: by replacing every connection $\circ - \circ$ with an arc, a directed graph results; we consider only *acyclic* versions.



The graphs all include two connections between V_1 and V_3 : the *simple* chain $V_1 \circ - \circ V_2 \circ - \circ V_3$ and the 'loopy' chain $V_1 \circ - \circ V_2 \circ - \circ V_4 \circ - \circ V_5 \circ - \circ V_2 \circ - \circ V_3$, which passes node V_2 twice. Recall that in the context of *blocking* and *d-separation* it is sufficient to consider *simple* chains only, where every node is included at most once.

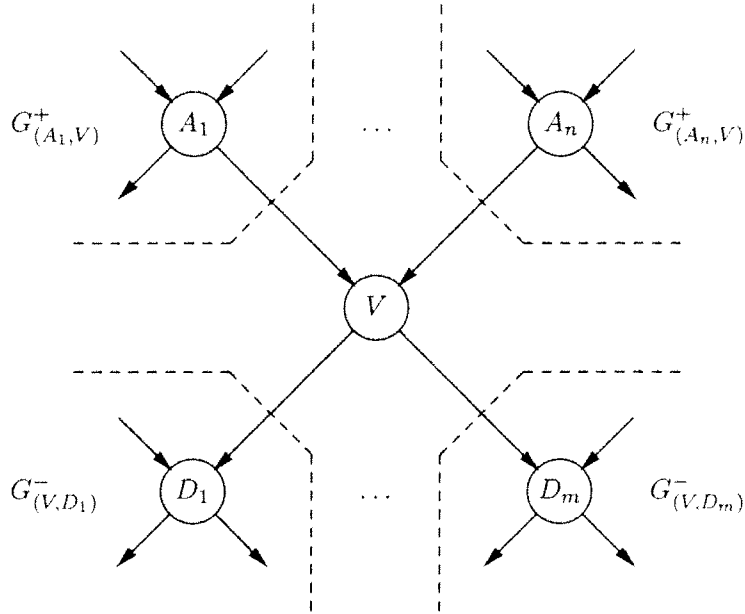
- Clearly demonstrate for all three graphs described above that the following holds regardless of the direction of the arcs for $V_2 \circ - \circ V_4$, $V_4 \circ - \circ V_5$ and $V_5 \circ - \circ V_2$:

if the simple chain between nodes V_1 and V_3 is blocked then the 'loopy' chain between these nodes is also blocked.

Hint: You may disregard all situations with evidence for V_1 or V_3 .

Formulas (Probabilistic Reasoning Exam Nov 2013)

Pearl in a singly connected digraph



Consider a node V in a probabilistic network $B = (G, \Gamma)$. Let $\rho(V) = \{A_1, \dots, A_n\}$ be the set of direct ancestors (parents) of V in G , and let $\sigma(V) = \{D_1, \dots, D_m\}$ be the set of its direct descendants (children). With Pearl's algorithm, node V computes the following parameters:

$$\pi(V) = \sum_{c_{\rho(V)}} \left(\gamma(V | c_{\rho(V)}) \cdot \prod_{i=1, \dots, n} \pi_V^{A_i}(c_{A_i}) \right)$$

$$\lambda(V) = \prod_{j=1, \dots, m} \lambda_{D_j}^V(V)$$

$$\pi_{D_j}^V(V) = \alpha \cdot \pi(V) \cdot \prod_{\substack{k=1, \dots, m \\ k \neq j}} \lambda_{D_k}^V(V)$$

$$\lambda_V^{A_i}(A_i) = \alpha \cdot \sum_{c_V} \lambda(c_V) \cdot \sum_{c_{\rho(V) \setminus \{A_i\}}} \left(\gamma(c_V | c_{\rho(V) \setminus \{A_i\}} \wedge A_i) \cdot \prod_{\substack{k=1, \dots, n \\ k \neq i}} \pi_V^{A_k}(c_{A_k}) \right)$$

in order to perform *data fusion*: $\alpha \cdot \pi(V) \cdot \lambda(V)$, which results in V 's (prior or posterior) probability distribution.

Problem 3 (a: 10 pts., b: 10 pts., c: 15 pts., total: 35 pts.)

Suppose that you are constructing a probabilistic network for a certain domain of application. For quantifying the network, a number of probabilities have to be assessed by a human expert. After full quantification, you will perform a sensitivity analysis to investigate the robustness of your network.

- a. Describe two different methods for eliciting probabilities from human experts. For each of the two methods, provide a benefit and a drawback.

Suppose the graph of your probabilistic network is the same as that from **Problem 1**.

- b. Give the *Sensitivity set* $S^{\{V_4, V_6\}}(V_3)$ for node of interest V_3 given evidence for nodes V_4 and V_6 . Clearly explain your answer.

Consider a sensitivity function $f(x)$ for some output probability and some parameter $x = \gamma_V(v_i)$ of variable V with $n > 2$ values. Upon varying x , the values of the assessment function γ_V for the other $n - 1$ values of variable V have to be co-varied. To this end, a *proportional* co-variation scheme is used in which the proportion of the remaining mass $1 - x$ that is assigned to $\gamma_V(v_j)$, $j \neq i$, is kept constant.

Different co-variation schemes adhere to different properties. Two useful properties are the *order-preserving* property and the *impossibility-preserving* property. A co-variation scheme is called impossibility-preserving if any $\gamma_V(v_j) = 0$, $j \neq i$, remains zero upon co-variation. Suppose the values of V are ordered according to the values of γ_V as specified in the network, i.e. $\gamma_V(v_1) \leq \dots \leq \gamma_V(v_n)$; a co-variation scheme is called order-preserving if this ordering is preserved during co-variation.

- c. Indicate which of the two properties described above are properties of the proportional co-variation scheme:
 - I. both order-preserving and impossibility-preserving
 - II. order-preserving, but not impossibility-preserving
 - III. not order-preserving, yet impossibility-preserving
 - IV. neither order-preserving, nor impossibility-preserving

For each of the properties provide a clear proof or counter-example.